

**STAT 309: MATHEMATICAL COMPUTATIONS I**  
**FALL 2013**  
**PROBLEM SET 5**

1. Consider an  $n \times n$  tridiagonal matrix of the form

$$T_\alpha = \begin{bmatrix} \alpha & -1 & & & \\ -1 & \alpha & -1 & & \\ & -1 & \alpha & -1 & \\ & & -1 & \alpha & -1 \\ & & & -1 & \alpha \end{bmatrix},$$

where  $\alpha \in \mathbb{R}$  is a real parameter.

(a) Verify that the eigenvalues of  $T_\alpha$  are given by

$$\lambda_j = \alpha - 2 \cos(j\theta), \quad j = 1, \dots, n,$$

where

$$\theta = \frac{\pi}{n+1}$$

and that an eigenvector associated with each  $\lambda_j$  is

$$\mathbf{q}_j = [\sin(j\theta), \sin(2j\theta), \dots, \sin(nj\theta)]^\top.$$

Under what condition on  $\alpha$  does this matrix become positive definite?

(b) If an iterative method with iteration matrix  $B_\alpha$  is convergent, we define its *asymptotic rate of convergence* as

$$\lim_{m \rightarrow \infty} \left( \frac{-\log \|B_\alpha^m\|_2}{m} \right)$$

Take  $\alpha = 2$ .

- (i) Will the Jacobi iteration converge for this matrix? If so, what will its asymptotic rate of convergence be?
- (ii) Will the Gauss-Seidel iteration converge for this matrix? If so, what will its asymptotic rate of convergence be?
- (iii) For which values of  $\omega$  will the SOR iteration converge?

2. In general, a *semi-iterative method* is one that comprises two steps:

$$\mathbf{x}^{(k+1)} = M\mathbf{x}^{(k)} + \mathbf{b} \tag{Iteration}$$

and

$$\mathbf{y}^{(m)} = \sum_{k=0}^m \alpha_k^{(m)} \mathbf{x}^{(k)}. \tag{Extrapolation}$$

As in the lectures, we will assume that  $M = I - A$  with  $\rho(M) < 1$  and that we are interested to solve  $A\mathbf{x} = \mathbf{b}$  for some nonsingular matrix  $A \in \mathbb{C}^{n \times n}$ . Let

$$\mathbf{e}^{(k)} = \mathbf{x}^{(k)} - \mathbf{x} \quad \text{and} \quad \boldsymbol{\varepsilon}^{(m)} = \mathbf{y}^{(m)} - \mathbf{x}.$$

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*Date:* November 26, 2013 (Version 1.2); due: 12 noon, December 9, 2013. Turn in your solutions to Kirsten Wellman in Eckhart 108. No late homeworks will be accepted.

- (a) By considering what happens when  $\mathbf{x}^{(0)} = \mathbf{x}$ , show that it is natural to impose

$$\sum_{k=0}^m \alpha_k^{(m)} = 1 \quad (2.1)$$

for all  $m \in \mathbb{N} \cup \{0\}$ . Henceforth, we will assume that (2.1) is satisfied for all problems in this problem set.

- (b) Show that for all  $m \in \mathbb{N}$ , we may write

$$\mathbf{e}^{(m)} = P_m(M)\mathbf{e}^{(0)}$$

for some  $P_m \in \mathbb{C}[x]$  with  $\deg(P_m) = m$  and  $P_m(1) = 1$ .

- (c) Hence deduce that a necessary condition for convergence is that

$$\lim_{m \rightarrow \infty} \|P_m(M)\|_2 < 1$$

where  $\|\cdot\|_2$  is the spectral norm. Is this condition also sufficient?

- (d) Consider the case when

$$\alpha_0^{(m)} = \alpha_1^{(m)} = \dots = \alpha_m^{(m)} = \frac{1}{m+1}$$

for all  $m \in \mathbb{N} \cup \{0\}$ . Show that if a sequence (any sequence, not necessarily one generated as in (Iteration)) is convergent and

$$\lim_{k \rightarrow \infty} \mathbf{x}^{(k)} = \mathbf{x}$$

then

$$\lim_{m \rightarrow \infty} \mathbf{y}^{(m)} = \mathbf{x}.$$

Is the converse also true?

- 3.** It is clear that in any semi-iterative method defined by some  $M \in \mathbb{C}^{n \times n}$  with  $\rho(M) < 1$ , we would like to solve the problem

$$\min_{P \in \mathbb{C}[x], \deg(P)=m, P(1)=1} \|P(M)\|_2. \quad (3.2)$$

Note that in the lectures, we required the polynomial  $P$  to satisfy  $P(0) = 1$ . Here we use a different condition,  $P(1) = 1$ , motivated by Problem 2(a).

- (a) Show that if  $m \geq n$ , then a solution to (3.2) is given by

$$P_m(x) = \frac{x^{m-n} \det(xI - M)}{\det(I - M)}.$$

How do we know that the denominator is non-zero?

- (b) From now on assume that  $M$  is Hermitian with minimum and maximum eigenvalues  $\lambda_{\min}, \lambda_{\max} \in \mathbb{R}$ . Define

$$\|f\|_{\infty} = \sup_{x \in [\lambda_{\min}, \lambda_{\max}]} |f(x)|.$$

Emulating our discussions in the lectures, show that for  $m = 0, 1, \dots, n-1$ , the solution to the relaxed problem

$$\min_{P \in \mathbb{C}[x], \deg(P)=m, P(1)=1} \|P\|_{\infty} \quad (3.3)$$

would yield an upper bound to (3.2).

- (c) Again by emulating our discussions in the lectures, show that the solution to (3.3) for  $\lambda_{\min} = -1$  and  $\lambda_{\max} = +1$  is given by the Chebyshev polynomials,

$$C_m(x) = \begin{cases} \cos(m \cos^{-1}(x)) & -1 \leq x \leq 1, \\ \cosh(m \cosh^{-1}(x)) & x > 1, \\ (-1)^m \cosh(m \cosh^{-1}(-x)) & x < -1. \end{cases}$$

- (d) Hence deduce that the solution to (3.3) for  $\lambda_{\min} = a$  and  $\lambda_{\max} = b$  is given

$$P_m(x) = \frac{C_m\left(\frac{2x - (b+a)}{b-a}\right)}{C_m\left(\frac{2 - (b+a)}{b-a}\right)}. \quad (3.4)$$

Note that this solves (3.3) for all  $m \in \mathbb{N}$  and not just  $m \leq n - 1$ .

- (e) Show that the solution in (d) is unique.

4. Let  $M \in \mathbb{C}^{n \times n}$  be Hermitian with  $\rho(M) = \rho < 1$ . Moreover, suppose that

$$\lambda_{\min} = -\rho, \quad \lambda_{\max} = \rho.$$

- (a) Show that the  $P_m$ 's in (3.4) satisfy a three-term recurrence relation

$$C_{m+1}\left(\frac{1}{\rho}\right) P_{m+1}(x) = \frac{2x}{\rho} C_m\left(\frac{1}{\rho}\right) P_m(x) - C_{m-1}\left(\frac{1}{\rho}\right) P_{m-1}(x)$$

for all  $m \in \mathbb{N}$ .

- (b) Show that the semi-iterative method with  $\alpha_k^{(m)}$  given by the coefficient of  $P_m$  in (3.4) may be written as

$$\mathbf{y}^{(m+1)} = \omega_{m+1}(M\mathbf{y}^{(m)} - \mathbf{y}^{(m-1)} + \mathbf{b}) + \mathbf{y}^{(m-1)}$$

where  $\omega_1 = 1$  and

$$\omega_{m+1} = \frac{2C_m(1/\rho)}{\rho C_{m+1}(1/\rho)}$$

for  $m = 0, 1, 2, \dots$ . This is a slightly different Chebyshev method where we choose the normalization (2.1) instead of  $\alpha_m^{(m)} = 1$  in the lecture.

- (c) Show that

$$\|P_m(M)\|_2 = \frac{1}{C_m(1/\rho)} = \frac{1}{\cosh(m\sigma)}$$

where  $\sigma = \cosh^{-1}(1/\rho)$ . Deduce that  $\|P_m(M)\|_2$  is a strictly decreasing sequence for all  $m = 0, 1, 2, \dots$

- (d) Show that

$$e^{-\sigma} = (\omega - 1)^{1/2}$$

where

$$\omega = \frac{2}{1 + \sqrt{1 - \rho^2}} \quad (4.5)$$

and deduce that

$$\|P_m(M)\|_2 = \frac{2(\omega - 1)^{m/2}}{1 + (\omega - 1)^m}.$$

- (e) Hence show that  $(\omega_m)_{m=0}^{\infty}$  is strictly decreasing for  $m \geq 2$  and that

$$\lim_{m \rightarrow \infty} \omega_m = \omega.$$

5. Let  $M \in \mathbb{C}^{n \times n}$  be nonsingular with  $\rho(M) < 1$  and suppose we are interested in solving

$$M\mathbf{x} = \mathbf{b}. \quad (5.6)$$

- (a) Show that SOR applied to the system

$$\begin{bmatrix} I & -M \\ -M & I \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ \mathbf{z} \end{bmatrix} = \begin{bmatrix} \mathbf{b} \\ \mathbf{b} \end{bmatrix} \quad (5.7)$$

yields the following iterations

$$\begin{aligned} \mathbf{x}^{(m+1)} &= \omega(M\mathbf{z}^{(m)} - \mathbf{x}^{(m)} + \mathbf{b}) + \mathbf{x}^{(m)}, \\ \mathbf{z}^{(m+1)} &= \omega(M\mathbf{x}^{(m+1)} - \mathbf{z}^{(m)} + \mathbf{b}) + \mathbf{z}^{(m)}, \end{aligned}$$

for  $m = 0, 1, 2, \dots$

- (b) Define the sequence of iterates  $\mathbf{y}^{(m)}$  by

$$\mathbf{y}^{(m)} = \begin{cases} \mathbf{x}^{(k)} & \text{if } m = 2k, \\ \mathbf{z}^{(k)} & \text{if } m = 2k + 1. \end{cases}$$

Show that the iterations obtained in (a) are exactly the iterations in Problem 4(b). This shows that SOR applied to (5.7) is equivalent to Chebyshev applied to (5.6) but with  $\omega_m = \omega$  for all  $m \in \mathbb{N}$ . Note that if  $\omega$  is chosen to be the value in (4.5), then this is in fact the optimal SOR parameter.

6. Let  $A \in \mathbb{R}^{n \times n}$  be symmetric positive definite and  $\mathbf{b} \in \mathbb{R}^n$ . As usual, we write

$$\mathbf{r}_k = \mathbf{b} - A\mathbf{x}_k. \quad (6.8)$$

We assume that  $\mathbf{x}_0$  is initialized in some manner. In the lectures we assumed  $\mathbf{x}_0 = \mathbf{0}$  and so  $\mathbf{r}_0 = \mathbf{b}$  but we will do it a little more generally here. Consider the quadratic functional

$$\varphi(\mathbf{x}) = \mathbf{x}^\top A\mathbf{x} - 2\mathbf{b}^\top \mathbf{x}.$$

- (a) Show that

$$\nabla \varphi(\mathbf{x}_k) = -2\mathbf{r}_k$$

and hence if  $\mathbf{x}_* \in \mathbb{R}^n$  is a stationary point of  $\varphi$ , then

$$A\mathbf{x}_* = \mathbf{b}.$$

Show also that  $\mathbf{x}_*$  must be a minimizer of  $\varphi$ .

- (b) Consider an iterative method

$$\mathbf{x}_{k+1} = \mathbf{x}_k + \alpha_k \mathbf{p}_k \quad (6.9)$$

where  $\mathbf{p}_0, \mathbf{p}_1, \mathbf{p}_2, \dots$  are search directions to be chosen later. Show that if we want  $\alpha_k$  so that the function  $f : \mathbb{R} \rightarrow \mathbb{R}$ ,

$$f(\alpha) = \varphi(\mathbf{x}_k + \alpha \mathbf{p}_k)$$

is minimized, then we must have

$$\alpha_k = \frac{\mathbf{r}_k^\top \mathbf{p}_k}{\mathbf{p}_k^\top A \mathbf{p}_k}. \quad (6.10)$$

- (c) Deduce that

$$\varphi(\mathbf{x}_{k+1}) - \varphi(\mathbf{x}_k) = -\frac{(\mathbf{r}_k^\top \mathbf{p}_k)^2}{\mathbf{p}_k^\top A \mathbf{p}_k}$$

and therefore  $\varphi(\mathbf{x}_{k+1}) < \varphi(\mathbf{x}_k)$  as long as  $\mathbf{r}_k^\top \mathbf{p}_k \neq 0$ .

(d) Show that if we choose

$$\mathbf{p}_k = \mathbf{r}_k, \tag{6.11}$$

we obtain the steepest decent method discussed in the lectures.

(e) Let the eigenvalues of  $A$  be  $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n > 0$  and  $P \in \mathbb{R}[t]$ . Show that

$$\|P(A)\mathbf{x}\|_A \leq \max_{1 \leq i \leq n} |P(\lambda_i)| \|\mathbf{x}\|_A$$

for every  $\mathbf{x} \in \mathbb{R}^n$ . [*Hint*:  $A \succ 0$  and so has an eigenbasis].

(f) Using (e) and  $P_\alpha(t) = 1 - \alpha t$ , show that if we have (6.11), then

$$\|\mathbf{x}_k - \mathbf{x}_*\|_A \leq \max_{1 \leq i \leq n} |P_\alpha(\lambda_i)| \|\mathbf{x}_{k-1} - \mathbf{x}_*\|_A$$

for all  $\alpha \in \mathbb{R}$ .

(g) Using properties of Chebyshev polynomials, show that

$$\min_{\alpha \in \mathbb{R}} \max_{\lambda_n \leq t \leq \lambda_1} |1 - \alpha t| = \frac{\lambda_1 - \lambda_n}{\lambda_1 + \lambda_n}$$

and hence deduce that

$$\|\mathbf{x}_k - \mathbf{x}_*\|_A \leq \frac{\lambda_1 - \lambda_n}{\lambda_1 + \lambda_n} \|\mathbf{x}_{k-1} - \mathbf{x}_*\|_A.$$

7. In this problem, you are required to randomly generate *sparse* matrices and sparse symmetric matrices of specified *densities*, i.e., proportion of nonzero elements, [cf. `sprandn` and `sprandsym` in Matlab/Octave/Scilab]. Modify the diagonal elements of your sparse matrices so that they are diagonally dominant and do likewise so that your sparse symmetric matrices are positive definite.

- (a) Implement Gaussian elimination with partial pivoting and Gauss-Seidel method. Compare their speeds and accuracies on a range of diagonally dominant matrices of varying dimensions and densities.
- (b) Implement Cholesky factorization and steepest descent method. Compare their speeds and accuracies on a range of symmetric positive definite matrices of varying dimensions and densities.
- (c) Comment on your findings in (a) and (b) with relevant numerical evidence presented in graphs and/or tables.