

# The Heisenberg Representation and the Fast Fourier Transform

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# Motivation:

- Discrete Fourier Transform

$$DFT = \frac{1}{\sqrt{N}} \left( e^{\frac{2\pi i}{N} \tau \cdot \omega} \right)_{0 \leq \tau, \omega \leq N-1}$$

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- Cooley–Tukey (1965):  $O(N \cdot \log(N))$  operations!

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$$L_\tau : \mathcal{H} \rightarrow \mathcal{H},$$
$$L_\tau[f](t) = f(t + \tau), \quad t \in \mathbb{Z}_N.$$



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- Note:

$$M_\omega \circ L_\tau = e^{\frac{2\pi i}{N}\omega\tau} \cdot L_\tau \circ M_\omega \quad \text{— Heisenberg commutation relations}$$

# Heisenberg Rep'n, cont.

- Combine:  $\tau, \omega, z \in \mathbb{Z}_N$

$$\pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau.$$

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- $(0, 0) \cdot (v, z) = (v, z) \cdot (0, 0) = (v, z);$

- $(v, z) \cdot (-v, -z) = (-v, -z) \cdot (v, z) = (0, 0).$

# Heisenberg Rep'n, cont.

- Summary: Heisenberg Rep'n

$$\left\{ \begin{array}{l} \pi : H \rightarrow GL(\mathcal{H}); \\ \pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau; \\ \pi(h \cdot h') = \pi(h) \circ \pi(h') \text{ — homomorphism.} \end{array} \right.$$

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A representation of a group  $H$  on a complex vector space  $\mathcal{H}$  is a homomorphism

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- Question. DFT ?

## (II) Representation Theory

### Definitions

We say that  $(\pi_1, H, \mathcal{H}_1)$ ,  $(\pi_2, H, \mathcal{H}_2)$  are equivalent,  $\pi_1 \simeq \pi_2$ , if

$$\exists \alpha : \mathcal{H}_1 \xrightarrow{\sim} \mathcal{H}_2 \text{ — Intertwiner,}$$

such that for every  $h \in H$

$$\begin{array}{ccc} \mathcal{H}_1 & \xrightarrow{\pi_1(h)} & \mathcal{H}_1 \\ \downarrow \alpha & & \downarrow \alpha \\ \mathcal{H}_2 & \xrightarrow{\pi_2(h)} & \mathcal{H}_2 \end{array}$$

i.e.,  $\alpha \circ \pi_1(h) = \pi_2(h) \circ \alpha$ .

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- Example: *DFT* is an intertwiner!

## Rep'n Theory, cont.

- $H = V \times Z = \underbrace{\mathbb{Z}_N}_T \times \underbrace{\mathbb{Z}_N}_W \times \underbrace{\mathbb{Z}_N}_Z$



# Rep'n Theory, cont.

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- Time representation:  $\mathcal{H}_T = \mathbb{C}(T)$

$$\begin{cases} \pi_T : H \rightarrow GL(\mathcal{H}_T); \\ \pi_T(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau. \end{cases}$$

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- Frequency representation:  $\mathcal{H}_W = \mathbb{C}(W)$

$$\begin{cases} \pi_W : H \rightarrow GL(\mathcal{H}_W); \\ \pi_W(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{-\frac{1}{2}\omega\tau + z\}} \cdot M_\tau \circ L_{-\omega}. \end{cases}$$

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- Fact:

$$DFT \circ \pi_T(h) = \pi_W(h) \circ DFT,$$

where

$$DFT[f](w) = \frac{1}{\sqrt{N}} \sum_{t \in \mathbb{Z}_N} f(t) e^{\frac{2\pi i}{N} wt}.$$

# (III) FFT Algorithm

- Idea:

$$\begin{array}{ccc} \mathcal{H}_T & \xrightarrow{\text{fast}} & \mathcal{H}^\Lambda \\ \text{slow} \downarrow \text{DFT} & & \parallel \\ \mathcal{H}_W & \xleftarrow{\text{fast}} & \mathcal{H}^\Lambda \end{array}$$

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- More representation theory:

## Definition

A rep'n  $(\pi, H, \mathcal{H})$  is irreducible if

$$\nexists 0 \neq \mathcal{H}' \subsetneq \mathcal{H}$$

such that

$$\pi(h) \cdot \mathcal{H}' \subset \mathcal{H}', \quad \forall h \in H.$$

## Theorem (Stone–von Neumann)

If  $(\pi_j, H, \mathcal{H}_j)$ ,  $j = 1, 2$ , are irreducible representations of the Heisenberg group  $H = V \times Z$ , with

$$\pi_j(z) = e^{\frac{2\pi i}{N}z} \cdot \text{Id}_{\mathcal{H}_j}, \quad \forall z \in Z,$$

then  $\pi_1 \simeq \pi_2$  are equivalent, i.e.,  $\exists \alpha : \mathcal{H}_1 \xrightarrow{\sim} \mathcal{H}_2$  such that

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## Theorem (Schur's lemma)

If  $\pi_1 \simeq \pi_2$  equivalent irreducible representations, and if  $\alpha, \alpha'$  satisfy equation (1), then

$$\alpha = c \cdot \alpha', \quad \text{for some scalar } c.$$



# FFT Algorithm: Models of Heisenberg Rep'n

## Examples

(1) Time model:  $V \supset T = \{(t, 0); t \in \mathbb{Z}_N\}$

$$\pi_T : H \rightarrow GL(\mathcal{H}_T).$$

(2) Frequency model:  $V \supset W = \{(0, w); w \in \mathbb{Z}_N\}$

$$\pi_W : H \rightarrow GL(\mathcal{H}_W).$$

## Corollary

*DFT* :  $\mathcal{H}_T \rightarrow \mathcal{H}_W$  — unique (up to a scalar) intertwiner between  $\pi_T$  and  $\pi_W$ .

## Examples

(3)  $W$ -invariant model:

Space:

$$\mathcal{H}^W = \left\{ \begin{array}{l} f : H \rightarrow \mathbb{C}, \\ f(w \cdot h) = f(h), \quad \forall w \in W, h \in H, \\ f(z \cdot h) = e^{\frac{2\pi i}{N}z} \cdot f(h), \quad \forall z \in Z, h \in H. \end{array} \right.$$

Action:

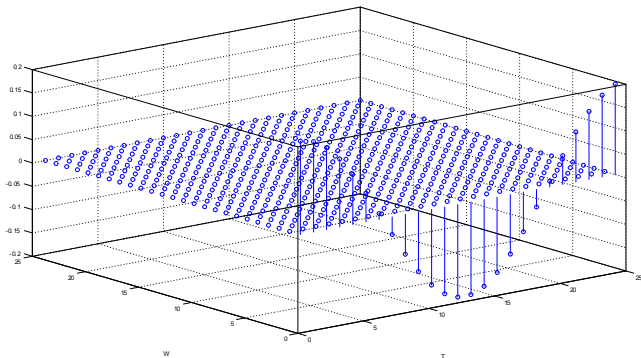
$$\left\{ \begin{array}{l} \pi^W : H \rightarrow GL(\mathcal{H}^W), \\ [\pi^W(h')f](h) = f(h \cdot h'). \end{array} \right.$$

- Remark: We have a natural identification  $\mathcal{H}_T = \mathcal{H}^W$ , given by

$$f \mapsto \tilde{f}(wtz) = e^{\frac{2\pi i}{N}z} \cdot f(t).$$

# Think on $W$ -invariant functions as functions on $T$

A function  $f(t)$  on  $T = \{(t, 0, 0); t \in \mathbb{Z}/5^2\}$



## Examples

(4)  $T$ -invariant model:

Space:

$$\mathcal{H}^T = \begin{cases} g : H \rightarrow \mathbb{C}, \\ g(t \cdot h) = g(h), \quad \forall t \in T, h \in H, \\ g(z \cdot h) = e^{\frac{2\pi i}{N}z} \cdot g(h), \quad \forall z \in Z, h \in H. \end{cases}$$

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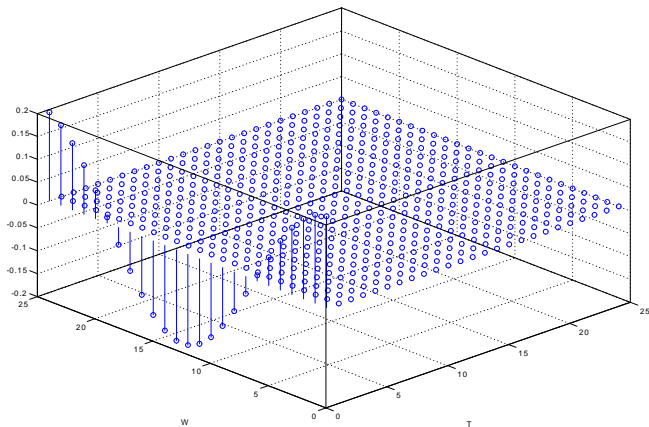
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# Think on T-invariant functions as functions on $W$

A function  $g(w)$  on  $W = \{(0, w, 0); w \in \mathbb{Z}/5^2\}$



# FFT Algorithm: The Arithmetic Model

## Examples

(5) Arithmetic model:  $N = p^2$

Lagrangian:  $V = \mathbb{Z}/p^2 \times \mathbb{Z}/p^2 \supset \Lambda = \{(p \cdot a, p \cdot b)\}$

Space:

$$\mathcal{H}^\Lambda = \begin{cases} F : H \rightarrow \mathbb{C}, \\ F(\lambda \cdot h) = F(h), \quad \forall \lambda \in \Lambda, h \in H, \\ F(z \cdot h) = e^{\frac{2\pi i}{N} z} \cdot F(h), \quad \forall z \in Z, h \in H. \end{cases}$$

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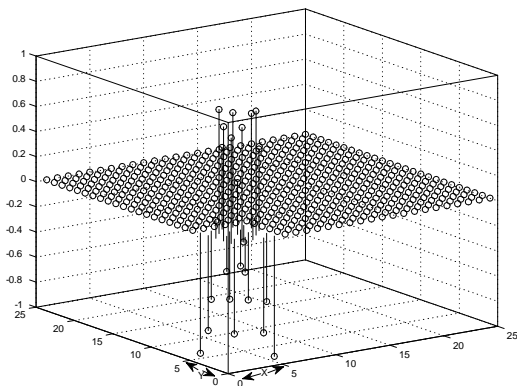
$$\pi^\Lambda : H \rightarrow GL(\mathcal{H}^\Lambda), \quad [\pi^\Lambda(h')F](h) = F(h \cdot h').$$

- Remark: We have a natural identification of  $\mathcal{H}^\Lambda$  with functions  $F(x, y)$  on  $\{0, \dots, p-1\} \times \{0, \dots, p-1\}$

$$F \mapsto \tilde{F}(\lambda \cdot (x, y, 0) \cdot z) = e^{\frac{2\pi i}{N} z} \cdot F(x, y), \quad \lambda \in \Lambda.$$

Think on Lambda-invariant functions as functions on  $\{0, \dots, p-1\} \times \{0, \dots, p-1\}$

A function  $F(x, y)$  on  $(x, y) \in \{0, \dots, 4\} \times \{0, \dots, 4\}$ ,  $p = 5$ ,



# FFT: The Algorithm

- Algorithm: Case  $N = p^2$

$$\begin{array}{ccc} \mathcal{H}_T & \xrightarrow[p^4]{DFT} & \mathcal{H}_W \\ \text{saw} \downarrow & & \uparrow \text{saw} \\ \mathcal{H}^W & & \mathcal{H}^T \\ \Sigma_{\lambda \in \Lambda} f(\lambda h) \downarrow ? & & ? \uparrow \Sigma_{t \in T} F(tw) \\ \mathcal{H}^\Lambda & \underline{\underline{=}} & \mathcal{H}^\Lambda \end{array}$$



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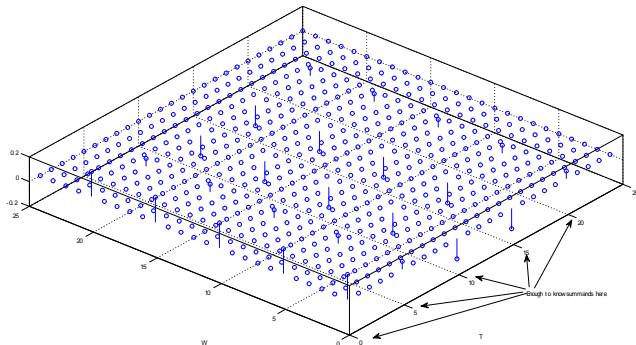
- Observation (**the saving!**): For  $f \in \mathcal{H}^W$ , if  $\lambda, \lambda' \in \Lambda$  differ by an element of  $W$ , i.e.,  $\lambda = w\lambda'$ , then

$$f(\lambda h) = f(w\lambda' h) = f(\lambda' h), \quad \forall h \in H.$$

# FFT Algorithms: The Summands

So for a fixed  $h$  it is enough to know the summands  $f(\lambda h)$ , in the transform, only for  $\lambda \in \Lambda/[\Lambda \cap W] = \{(0, 0), (p, 0), \dots, ((p-1)p, 0)\}$ .

- Example for  $p = 5$



# FFT Algorithm: The Saving

- So we have

$$\sum_{\lambda \in \Lambda} f(\lambda h) = \sum_{\lambda \in \Lambda / [\Lambda \cap W]} f(\lambda h) \cdot \#(\Lambda \cap W).$$

Only  $p = p^2 / p = \#(\Lambda / [\Lambda \cap W])$  summands !!!.

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- Complexity of the algorithm

$$\underbrace{p^2}_{\text{values of } h} \cdot \underbrace{p}_{\text{summands}} + \underbrace{p^2 \cdot p}_{\text{second operator}} = p \cdot p^2 \cdot (1 + 1) = \underbrace{p}_{\text{constant}} \cdot N \cdot \log(N).$$

# FFT Algorithm: Schur's lemma

Denote the intertwiners by

$$FFT^{\Lambda, W} : \mathcal{H}^W \rightarrow \mathcal{H}^{\Lambda}, \text{ and } FFT^{T, \Lambda} : \mathcal{H}^{\Lambda} \rightarrow \mathcal{H}^T.$$

- Why

$$DFT = FFT^{T, \Lambda} \circ FFT^{\Lambda, W} ? \quad (2)$$

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- Why

$$DFT = FFT^{T, \Lambda} \circ FFT^{\Lambda, W} ? \quad (2)$$

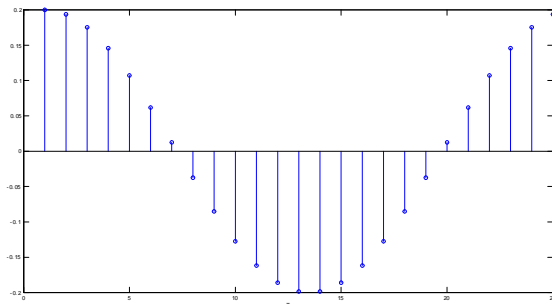
- In fact by Schur's lemma: Both sides of (2) intertwine  $\pi^W$  and  $\pi^T$ , so

$$DFT = c \cdot FFT^{T, \Lambda} \circ FFT^{\Lambda, W}$$

for some scalar  $c$ . Easy to compute  $c = \frac{1}{p}$ .

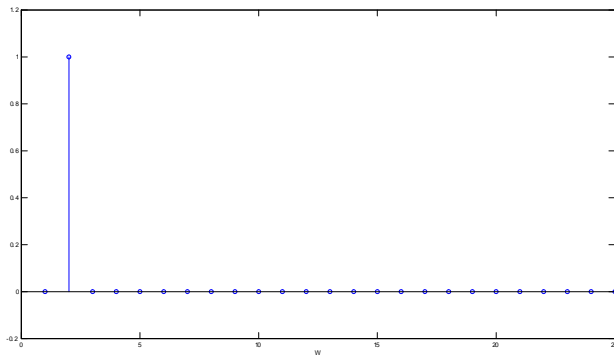
# Numerics for $n=25$

Start with the exponent function  $f(t) = e^{\frac{2\pi i}{25}t}/5$  on  $T = \mathbb{Z}/25$



# Numerics for $n=25$ : $\text{FFT}(f)$

Apply  $\text{FFT}$  and get  $g(w) = \text{FFT}[f](w)$  on  $W = \mathbb{Z}/25$



Indeed this is  $g(w) = \delta_1(w)$ .



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- FFT Algorithm

$$\begin{aligned} \mathcal{H}^W \xrightarrow{FFT^{\Lambda_1, W}} \mathcal{H}^{\Lambda_1} &\rightarrow \dots \rightarrow \mathcal{H}^{\Lambda_{k-1}} \xrightarrow{FFT^{T, \Lambda_k}} \mathcal{H}^T \\ DFT &= \frac{1}{\sqrt{p^k}} \cdot FFT^{T, \Lambda_{k-1}} \circ \dots \circ FFT^{\Lambda_1, W}, \end{aligned}$$

with complexity

$$\underbrace{p^k}_{\text{values of } h} \cdot \underbrace{\{k\}}_{k \text{ operators}} \cdot \underbrace{\# \Lambda_{j+1} / (\Lambda_j \cap \Lambda_{j+1})}_{p \text{ summands}} = \underbrace{p}_{\text{constant}} \cdot \underbrace{p^k}_N \cdot \underbrace{\log(p^k)}_{\log(N)}.$$

- Conclusion:

$\Lambda_{\bullet} \implies$  Cooley–Tukey *FFT*

Thank You!